LECTURE 1 INTRODUCTION

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408450 Computer Networks, Fall 2011/2012 http://www.hlms.hu.edu/

Course Staff

□ Instructors

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- □ Watch the course web page.
 - Handouts, readings, ..
- Read course boards.
 - "Announce" for official announcements
 - General" for questions/answers
- Books

Contact information

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Today's Lecture

- □ Course outline and goals.
- □ Tour of Networking

Contact Information II

Office Hours

Is posted at my office

D IM

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□ IM Hours

Variable

Note on Textbook

- Data Communications and Networking Behrouz a.
 Forouzan 4th Edition , 2007.
- We have to derive some material from other sources
 - Rapidly changing field
 - Most books focus on one or two specific topics in great detail
- □ I may not stick to the flow of the textbook

Course Policies

□ Your work MUST be your own

- No copying from web or other books without understanding the material
- Zero tolerance for cheating: You get an F for the course if you cheat in anything however small ² NO DISCUSSION
- □ Homework is due a week after it is assigned
 - Late assignments will NOT be accepted
- □ Check the webpage for everything!
 - You are responsible for checking the webpage for updates

Grading

- 25% for Exam I
- 25% for Exam II
- □ 50% for Final exam
- □ 0% for Homeworks & Quizes

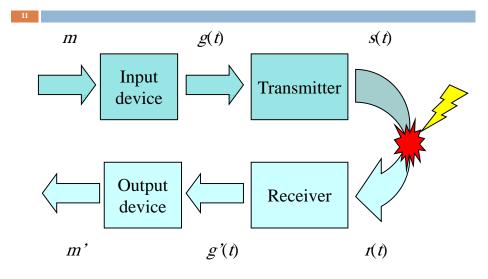
Exam Policies

- Writing must be clear, to the point, legible and unambiguous
- I should be able to understand what you are trying to do/say without verbal explanations later
- No credit for vague answers, unclear steps, magical solutions, etc.

Communication models

- The fundamental purpose of data communication is to exchange information between two entities
 - Information is represented as data and carries meaning currently assigned (depending on a context) to those data
 - Data represents facts, concepts, or instructions in a formalized manner suitable for communication, or processing by human or machines

Simple Communication Model



Simple Communication Model

- m: original information message created by sending entity
- \Box g(t): time varying signal, not suitable for transmission
- s(t): signal obtained by converting g(t) into a form that matches the characteristics of the transmission medium
- \Box r(t): received signal which may be different from s(t)
- g'(t): signal obtained by converting r(t) into a form suitable for output
- \square m': estimated message produced by destination entity

Simple Communication Model

- Although simple, the model highlights several communication tasks:
 - Interfacing,
 - Signal generation,
 - Synchronization,
 - Error detection and correction,
 - Flow control,
 - Framing, message formatting

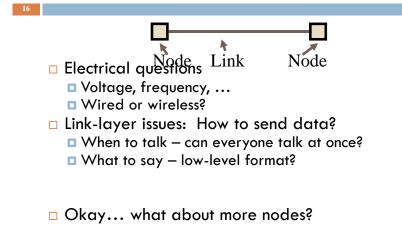
How to Draw a Network

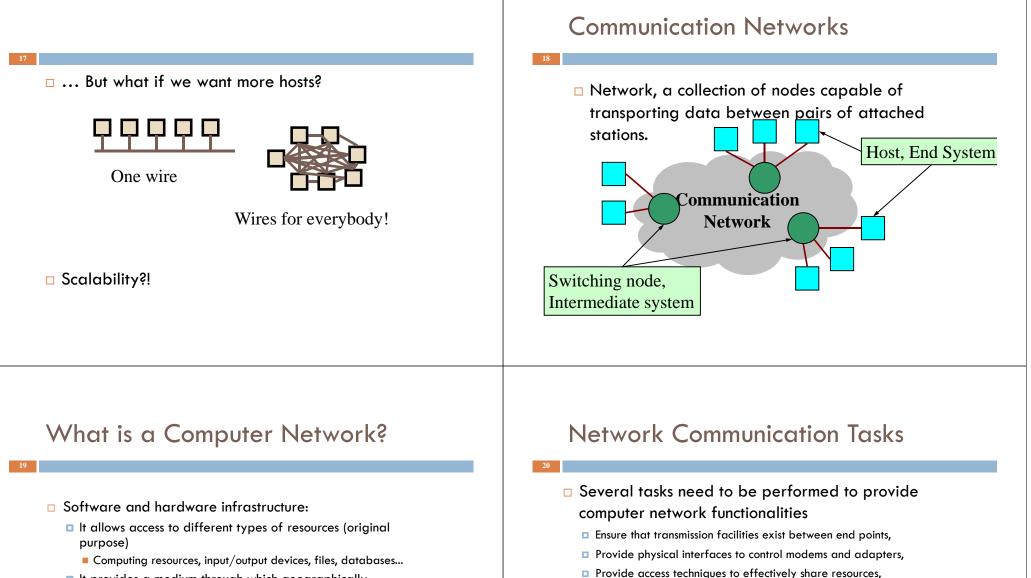


Data Communication Networks

- Point-to-point communication is the simplest form of data communication
 - Not practical to provide a dedicated wire between each pair of devices
- Communication networks provide efficient means for data communication between multiple devices

Building block: The Links





Provide addressing, routing, buffering and flow regulation of

Provide management functions and mechanisms to support

packets

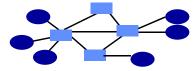
Accomodate peculiarities of users

distributed applications

- It provides a medium through which geographically dispersed users communicated (e.g. e-mail, chatting, teleconferencing)
- An electronic village
- An information highway, national information infrastructure
- Cyberspace "A consensual [environment] experienced daily by billions of operators, in every nation, ..."

Multiplexing!

Need to share network resources



- □ How? Switched network
 - Party "A" gets resources sometimes
 - Party "B" gets them sometimes
- □ Interior nodes act as "Routers" or "Switches"
- □ What mechanisms can share resources?

Classes of Networks

- □ Networks are usually defined according to
 - Geographical extent
 - Purpose
 - Implementation

Classes of Networks

- Based on geographical extent, the following classes of networks are identified:
 - Local Area Networks (LANs)
 - Metropolitan Area Networks (MANs)
 - Wide Area Networks (WANs)
 - Radio and Satellite Networks

Local Area Networks (LANs)

- □ By far, most commonly used class of networks.
 - Maximum distance between nodes is limited to few kilometers.
 - Usually owned by the same organization.
 - Typical data transmission ranges between 10 Mbps (coax cable) to 100 Mbps (fiber optic, category 5 unshielded twisted pair),
 - Several architectures have been standardized (802 LANs):
 - Ethernet (CSMA/CD), Fast Ethernet
 - Token Bus
 - Token Ring
 - FDDI

Metropolitan Area Networks (MANs)

- A relatively new class of networks, intermediate between LANs and WANs,
- Support moderate (Data, Voice and Video at 1 Mbps) to high data rates,
- MANs are optimized for a larger geographical area than LANs, ranging from several blocks of buildings to entire cities,
- Standards include FDDI, FDDI-II and IEEE 802.6 (DQDB).

Wide Area Networks (WANs)

- □ Cover a wide geographical area,
- Communication links are provided by telephone companies, or other common carriers,
- □ Transmission speeds are usually restricted,
- □ Transmission quality is inferior to LAN one,
- WAN's architechture is usually more complex than LAN's:
 - Need for efficient routing and congestion control algorithms,
- □ Many implentations of vendor architectures.

Radio and Satellite Networks

- □ Two basic subclasses of wireless networks:
 - Terrestial radio networks,
 - Satellite networks.

Communication Networks

- Communication networks can be classified as:
 - Broadcast networks, or
 - Switched networks.

Broadcast networks

- In any broadcast network, stations communicate over a shared medium,
- □ No need for intermediate node,
- □ Need for medium access control algorithm.
- □ Three types of broadcast networks:
 - Packet radio networks,
 - Satellite networks, and
 - Local area networks.

Switched networks

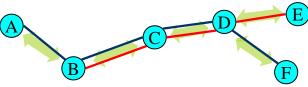
- - In switched networks, communication facilities are shared among users.
 - □ Three primary switching techniques:
 - Circuit switching
 - Message switching, and
 - Packet switching.

Switching Techniques: Circuit Switching

- □ Commonly used in the public telephone system.
- Exclusive dedication of a portion of the available bandwidth to carry traffic between a source and a destination.
- Allocation of the required bandwidth is achieved using:
 - Frequency Division Multiplexing (FDM)
 - Time Division Multiplexing (TDM)

Circuit Switched Networks

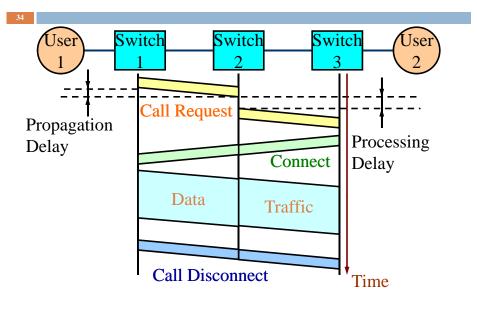
- Telephone network typically
 - All resources (e.g. communication links, buffers) needed by a "call" are reserved for its entire duration,
 - Resource reservation (i.e. resources are always available when needed by a call) guarantees "quality of service" (QoS)



Circuit Switching

- □ Call requires three phases:
 - Connection phase: a circuit is set up between source and destination,
 - Transmission phase: traffic exchange takes place,
 - Termination phase: the call is disconnected.

Circuit Switching Networks

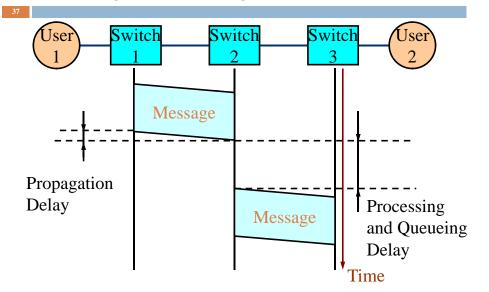


Circuit Switching Networks

- Delays for setting up connections can be high.
 - Ordinary telephone lines:
 - Call setup is on the order of 5 to 25 seconds after completion of dialing.

Message Switching Networks

- □ A physical circuit is shared among multiple users.
 - Leased communication facilities are used.
- Data enters the network in the form of "messages"
 - Messages are stored and subsequently forwarded.
 - No circuit switching delays are involved.
 - Queueing delays occur.
 - Message lengths are slightly longer because of headers.

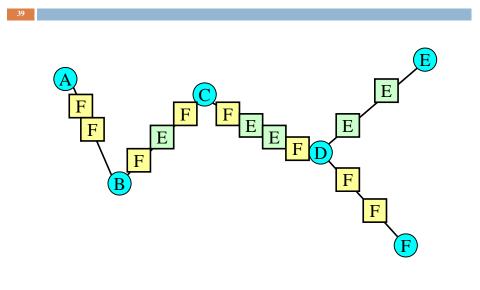


Message Switching Networks

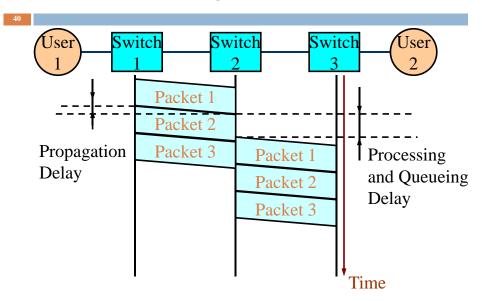
Packet Switching Networks

- Equivalent to message switching for short messages.
 - Maximum message length for transmission is imposed.
 - Any message exceeding the maximum is broken up into shorter units called "packets".
- Packets traversing a network share netwroks resources with other packets – statistical sharing of resources
 - Demand for resources may exceed amount of the resources available: contention

Packet Switching Networks



Packet Switching Networks



Performance Tradeoffs

□ Three types of delay:

- Propagation delay: time for the signal to propagate from the source to the receiver (depends on the wave propagation speed, in a wire 200 mln m/s = 200 m/µs),
- Transmission time: time needed to transmit the signal representing a block of data – depends on a link data rate,
- Processing and queueing delay: time needed to perform tasks necessary to relay a message/packet from one link to another in a given node, as well as time a message/packet has to spend waiting for the access to transmission medium – depends on the processing speed and the network load.

Performance Tradeoffs – Packet Switching

- □ A major benefit: the pipelining effect.
- □ Simultaneous use of communication circuits allows:
 - Considerable gain in efficiency,
 - Shorter delays, despite inclusion of headers for each packet.
 - Lower probability of retransmission,
 - Shorter messages are less likely to have errors than longer ones,
 - Errors do not cause retransmission of entire messages, but only of relatively shorter packets.
 - Packets can be routed independently, possibly minimizing congestion.

Performance Tradeoffs

- Packet switching provides flexibility in meeting the user needs.
 - Example: the needed rate is 75 kbps while the channel rate is 64 kbps
 - The use of packet switching meets the demands of the user more easily.

Performance Tradeoffs

- Statistical sharing of resources can be more efficient
 - Consider a 1 Mbps communication link
 - □ Each user requires 100 kbps when transmitting, but sends 10% of the time,
- □ Circuit switching:
 - Each caller is allocated 100 kbps capacity,
 - At most 10 callers are supported.
- Packet switching:
 - With 35 ongoing calls, probability that 10 or more callers are simultaneously active is about 0.00174,
 - Can support many more callers, with small probability of contension
- If user traffic is "bursty" (on/off), then packet switching can be more efficient than circuit switching.

Performance Tradeoffs

- The relative performance of circuit switching and packet switching depends on:
 - $\hfill\square$ The quality of service requirements of the application
 - End-to-end delay bounds,
 - Jitter control.
 - □ The traffic pattern,
 - Burtiness vs constant bit rate,
 - Connection setup overhead.

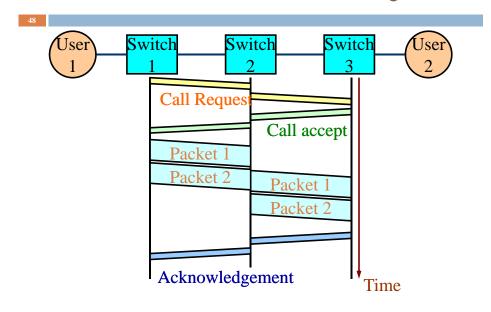
Packet Switching Techiques

- Two basic approaches to packet switching are common:
 - Virtual circuit packet switching
 - Datagram packet switching

Virtual Circuit Packet Switching

- □ An initial phase is used to setup a fixed route.
 - Similar to circuit switching, except that a delay occurs at each node,
 - Call request and call accept must both wait their turns on transmission.

Virtual Circuit Packet Switching



Virtual Circuit Packet Switching

- Upon path setup, the virtual circuit appears to the user as a dedicated circuit.
 - In reality, the circuit is shared among multiple users.
- Destination address is no longer required.
 - Only a virtual circuit number is needed to identify the destination.
 - Packets have shorter headers and fixed routing makes fast packet switching possible.

Datagram Packet Switching

- Datagram Packet Switching does not require a call setup
- □ For short transactions, it may be faster
- Individual datagrams are routed independently
 - Increases processing overhead at the router
 - Routing table lookups

Virtual Circuit Service Characteristics

- □ Guaranteed, reliable delivery
 - Powerful error control
 - Sequencing of packets
 - Detection and suppression of duplicates
- Congestion control minimizes queueing delays
 - Delays, however, are more variable than they are with dedicated circuits
- □ Enhanced security

Datagram Service Characteristics

- The network makes a "best effort" attempt to deliver the packets
 - Each packet is treated as a separate entity with no prior route determination
 - Packets may follow different paths to destination
 - No guarantees for reliable delivery
 - Packets may be lost, duplicated, or may arrive out of order
- The network relies on the user application to enhance the basic datagram service

Analogy

- □ Telephone service ⇔ virtual circuit
 - User must set up the connection, transmit and finally disconnect.
 - End users have the illusion that they communicate through a dedicated circuit.
 - Data are received in the order they have been transmitted.
- Regular postal service datagram
 - Letters are handled independently.
 - No guaranteed delivery. Losses are the user's responsibility.
 - Letters are not necessarily delivered in the order they have been sent.

Analogy

- □ User perspective:
 - " "Unsophisticated" users prefer virtual circuits.
 - The communication subnet handles most of the functions required to provide the requested service.
 - "Sophisticated" users prefer datagram service.
 - Datagram service provides more flexibility to implement particular features.

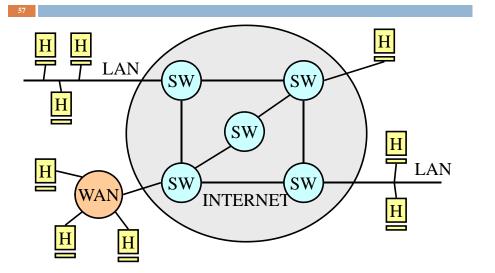
Analogy

- Application perspective:
 - Prior establishment of connection, although desirable, may lead to unacceptable delays for certain types of applications.
 - Request/Reply
 - Features such as error control may be detrimental to time bound applications, such as voice.
 - Few bits in error is far more preferable than lengthly retransmission delays.
 - For voice application a connection set up may be needed but no error control is required.

Connection Semantics

- Literature often uses the term connection-oriented and connectionless to refer to different network services.
 - Virtual circuit transmission is a special case of connection-oriented transmission
 - Datagram service is a special case of connectionless transmission

Internetworking



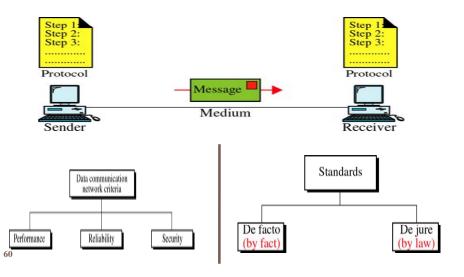
Internetworking

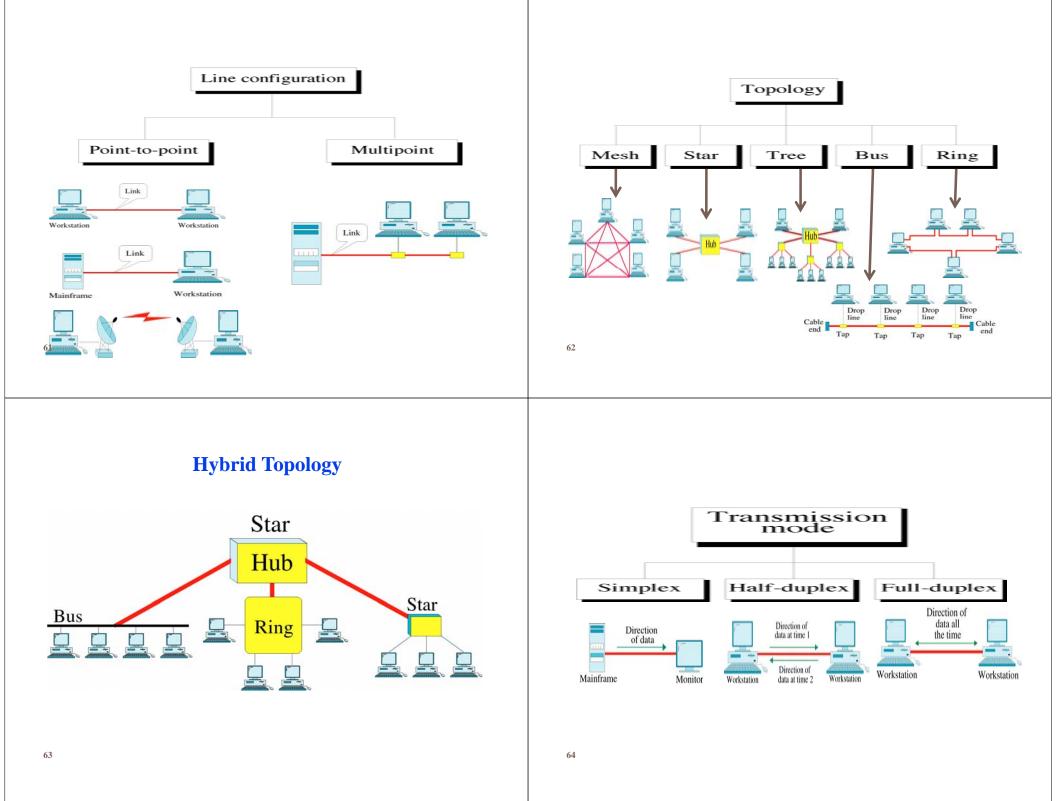
- The main goal is to provide a universal network formed out of physically different networks.
- □ Internetworking involves complex issues:
 - Different addressing and naming schemes
 - Different routing techniques
 - Different congestion control techniques
 - Different hardware interfaces
 - Connection oriented vs connectionless services
 - Different data unit sizes
 - Different error control techniques

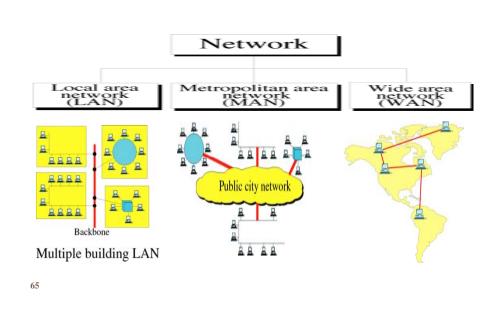
Summary

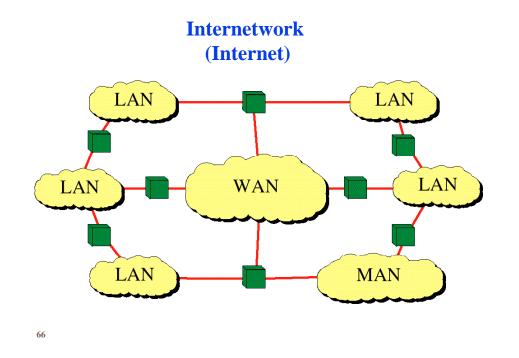
- Introduced the basic concepts of communication networks
- Discussed switching techniques
- Presented different classes of communication networks
- □ Introduced the concept of Internetworking

Data Communication System Components









"The Internet"

- □ An inter-net: a network of networks.
 - A set of networks that are connected with each other
 - Networks are connected using routers that support communication in a hierarchical fashion
 - Often need other special devices at the boundaries for security, accounting, ..
- The Internet: the interconnected set of networks of the Internet Service Providers (ISPs) providing data communications services.
 - About 17,000 (much more !!!) different networks make up the Internet
- In order to inter-operate, all participating networks have to follow a common set of rules.

Challenges of the Internet

- □ Scale: 100,000,000s of hosts
- □ Heterogeneity:
 - 18,000+ administrative domains
 - Thousands of applications
 - Lots of users
 - Fast links, slow links, satellite links, cellular links, carrier pigeons
- Diversity of network technologies
- Adversarial environment
- □ Oh, and let's make it easy to use...

Implementing Packet-Switched

Networks

- Requirements for packets:
 - Header information: Addresses, etc.
 - Data. What is packet size limit?
 - Everybody has to agree on these for interoperability
- □ How do packets reach destination? Routing
 - Nodes in network forward packets towards destination
 - Routing tells nodes where to send the packets they receive
 - Design questions: What criteria to decide?
 - Destination is a must
 - Source?
 - "Type"?

Routing

- □ Who chooses the routes?
 - A human: Static routing
 - Centralized routing (telenet, c.a. 1980s)
 - Distributed routing (Internet, ...)
- Distributed routing uses a Routing Protocol
 - Many different protocols are in use.
 - □ Inside an organization: RIP, OSPF, etc
 - Between organizations: BGP

Network Security

The field of network security is about:

- how bad guys can attack computer networks
- how we can defend networks against attacks
- how to design architectures that are immune to attacks
- Internet not originally designed with (much) security in mind
 - original vision: "a group of mutually trusting users attached to a transparent network" ⁽²⁾
 - Internet protocol designers playing "catch-up"
 - Security considerations in all layers!

Bad guys can put malware into hosts via Internet

- □ Malware can get in host from a virus, worm, or trojan horse.
- Spyware malware can record keystrokes, web sites visited, upload info to collection site.
- Infected host can be enrolled in a botnet, used for spam and DDoS attacks.
- Malware is often self-replicating: from an infected host, seeks entry into other hosts

Bad guys can put malware into hosts via Internet

Trojan horse

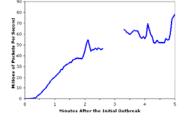
- Hidden part of some otherwise useful software
- Today often on a Web page (Active-X, plugin)

\Box Virus

- infection by receiving object (e.g., e-mail attachment), actively executing
- self-replicating: propagate itself to other hosts, users

□ Worm:

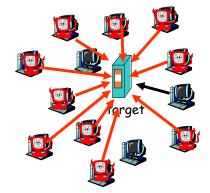
- infection by passively receiving object that gets itself executed
- self- replicating: propagates to other hosts, users
- Sapphire Worm: aggregate scans/sec in first 5 minutes of outbreak (CAIDA, UWisc data)



Bad guys can attack servers and network infrastructure

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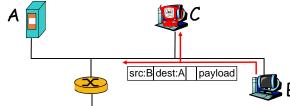
- Denial of service (DoS): attackers make resources (server, bandwidth) unavailable to legitimate traffic by overwhelming resource with bogus traffic
- 1. select target
- break into hosts around the network (see botnet)
- send packets toward target from compromised hosts



The bad guys can sniff packets

Packet sniffing:

- broadcast media (shared Ethernet, wireless)
- promiscuous network interface reads/records all packets (e.g., including passwords!) passing by



 Wireshark software used for end-ofchapter labs is a (free) packet-sniffer

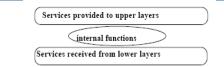
Layering Concept

- The complete communication system is broken down to a set of layers. Each layer will serve the one above it, and uses the services provided by the layer below it.
- The layering strategy is adopted in almost all organizations in order to simplify, and streamline the operation.

Definitions:

- Protocol: a certain set of rules that must be followed by all the systems that wish to communicate
- Peer processes: similar processes on different machines/systems/networks doing the same job & communicate according to a certain protocol
- Network architecture: is the set of network layers

Layer Interaction



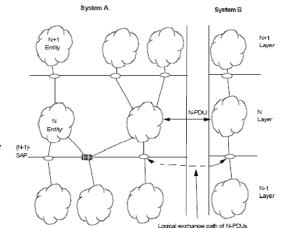
- □ Each layer (N) consists of one, or more, N-entities (may be different).
- □ N-entities in one system communicate with N-entities in another system.
- Entities are active processes that implement the layer's functions.
- N-Entities in different systems exchange N-layer protocol data units (N-PDU) (virtually)
- N-entities provide services to (N+1)-entities, and receive service from (N-1) entities:
 - N-Protocol Data Unit (N-PDU) is passed with Interface Control Information (ICI) to N-1 layer: Interface Data Unit (IDU)
 - (N-1)-Service Data Unit (N-SDU) is received and processed by layer N-1

Layer Interaction

Service primitives specify the type of service

• *Service parameters* are passed to service primitives (N-entities contain several service primitives)

• Interface to N-entities is through *service access points* (SAP).



Layer Design Issues

□ Connection mode

- Connection-oriented
- Connectionless
- Data Unit Size: Maximum and minimum data unit sizes
 - maximum data unit size is required for error detection, and to prevent hogging of the channel
 - minimum data unit size is required for efficient utilization of the bandwidth, and for the correct operation of some protocols
- □ If a data unit is too large:
 - use segmentation at the source: break the data unit into shorter segments
 - use reassembly at the destination: combine the shorter segments into the original data unit
- □ If a data unit is too small:
 - use blocking at the source: combine several data units into one
 - use unblocking at the destination: retrieve the original PDUs

- Data transfer mode (types of channels)
 - Simplex, Half Duplex, Full Duplex
- □ Addressing: source and destination addresses
- Logical channel creation and identification

The OSI Reference Model

The Seven Layer Model:

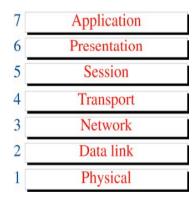
- The OSI reference model consists of 7 layers
- The layers are divided into 3 enclosed sets (environments):

1. Network Environment: Concerned with communication over the network.

2. OSI Environment :

Network environment layers + layers that serve the application.

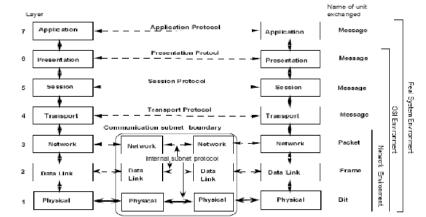
- 3. Real System Environment:
- OSI Environment + application itself.



OSI Model



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Network Environment

1. The Physical Layer: Physical & electrical interface to the communication channel.

Physical: Medium of transmission (wires)

Electrical: Current, voltage or optical sources, Modulation technique, Establishment & termination of connections

2. The Data Link Layer: Provides error-free data transmission between individual machines on the same network: Error checking; Request for retransmission; Acknowledgments; Detection of duplicates; Flow (congestion) control.

3. The Network Layer: Responsible for packet delivery network-wide: Routing across the network, Addressing, Reliable network-wide data delivery, Flow control (Across network)

The above 3 layers serve *network-related* functions: Deliver data safely, and in order between any 2 machines

OSI Environment

4. The Transport Layer: Provides reliable transfer of data independent of the type and number of networks involved:

- Breaks long messages into packets @ transmitter
- Creates messages from packets @ receiver
- Implement broadcast and multicast services
- Multiplex several transport connections on one network connection
- May create several network connections in order to serve one transport connection

• Provides several (5) classes of service, depending on network characteristics. **5.** The Session Layer: It manages sessions by providing primitives that handle

the station dialogs. It can be absent, or merged with the transport layer. **6.** The Presentation Layer: Provide machine-independent data representation:

Performs data conversion to and from an abstract form, in order to overcome code incompatibilities

Real System Layers

7. The Application Layer: User Applications, e.g., electronic mail & file transfer, Factory-floor communication, Distributed data-base systems, Network file servers, Network virtual terminal

Example:	Layer	Function
	Application	File Transfe
	Presentation	Convert to t
	Session	Set up & ter
	Transport	Break file ir
	Network	Route packe
	Data Link	Transmit pa

Physical

	Function
	File Transfer
ı	Convert to the abstract notation
	Set up & terminate connection
	Break file into packets in sequence
	Route packets to destination
	Transmit packets error-free between
	adjacent nodes; flow control
	Transmit raw data on cable

Layer Functions

	Application	To allow access to network resources
To translate, encrypt, and compress data	Presentation	To establish, manage,
To provide end-to-end	Session	and terminate
message delivery and	Transport	sessions To move packets from
To organize bits into	Network	source to destination; to provide internetworking
frames; to provide node- to-node delivery	Data link	To transmit bits; to provide
to-node derivery	Physical	mechanical and electrical specifications

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The TCP/IP Reference Model

 The TCP/IP suite of protocols was developed during the ARPANET project.

1. The physical layer:

connectionless protocol

defines the interface between the physical medium and the communicating device, e.g., twisted pairs, fibers, etc.
2. The network access layer: defines procedures for data exchange between the communicating system and the network, e.g., Ethernet, token ring, etc
3. The internet layer: defines procedures for routing and addressing of data

Units, e.g., The Internet Protocol (IP) is used for addressing, and routing, and is a

5. Application layer
4. Transport layer
3. Internet layer
2. Network access layer
1. Physical layer

Application

₽

Ethernet

Device driver and

hardware

UDP

TCP