

Data Communications and Networking Fourth Edition

### Forouzan

# **Chapter 21**

# Network Layer: Address Mapping, Error Reporting, and Multicasting

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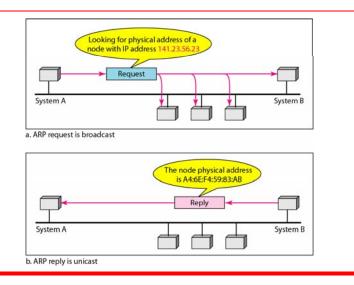
### **21-1 ADDRESS MAPPING**

The delivery of a packet to a host or a router requires two levels of addressing: logical and physical. We need to be able to map a logical address to its corresponding physical address and vice versa. This can be done by using either static or dynamic mapping.

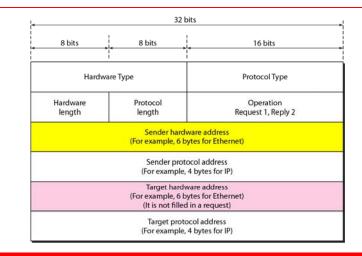
**Topics discussed in this section:** Mapping Logical to Physical Address Mapping Physical to Logical Address

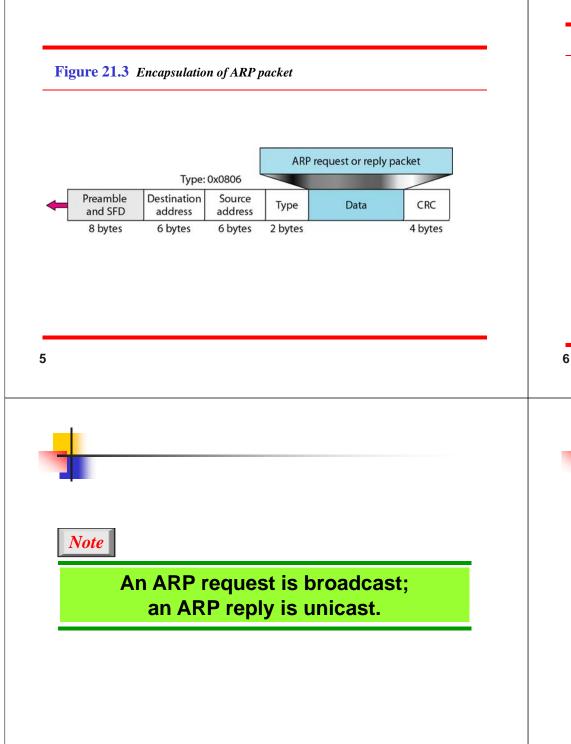
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#### Figure 21.1 ARP operation

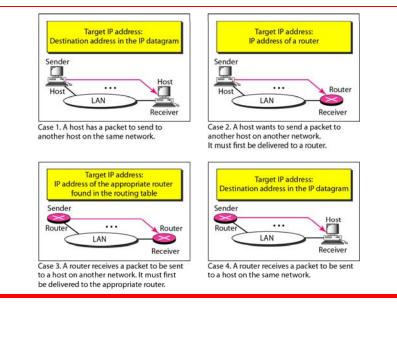


#### Figure 21.2 ARP packet





#### Figure 21.4 Four cases using ARP

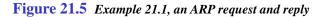


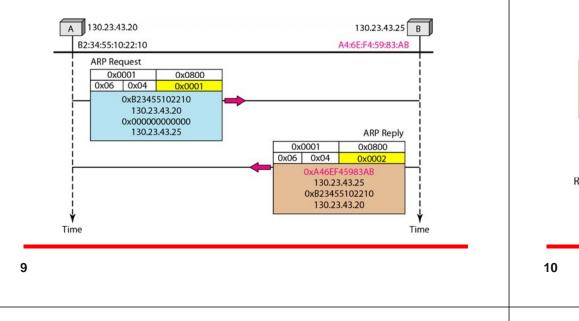
### Example 21.1

A host with IP address 130.23.43.20 and physical address B2:34:55:10:22:10 has a packet to send to another host with IP address 130.23.43.25 and physical address A4:6E:F4:59:83:AB. The two hosts are on the same Ethernet network. Show the ARP request and reply packets encapsulated in Ethernet frames.

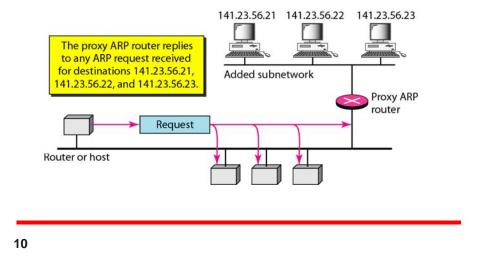
#### **Solution**

Figure 21.5 shows the ARP request and reply packets. Note that the ARP data field in this case is 28 bytes, and that the individual addresses do not fit in the 4-byte boundary. That is why we do not show the regular 4-byte boundaries for these addresses.





#### Figure 21.6 Proxy ARP



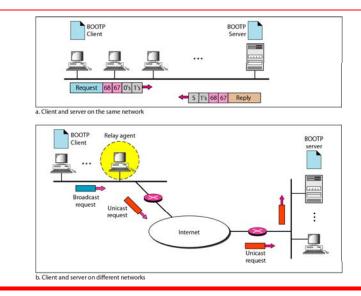
### Mapping Physical to Logical Address

When a station is just booted, or an organization does not have enough addresses to assign to each station, it will do it on demand.

- RARP (Broadcasting is done at the data link layer)
- BOOTP (Application layer processes)



Figure 21.7 BOOTP client and server on the same and different networks

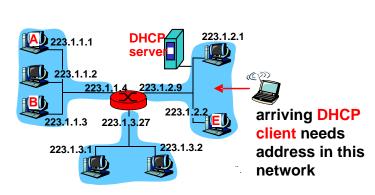


#### Note

**DHCP** provides static and dynamic address allocation that can be manual or automatic.

#### 13

### **DHCP** client-server scenario



### **DHCP: Dynamic Host Configuration Protocol**

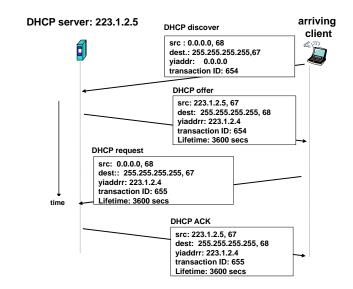
Goal: allow host to dynamically obtain its IP address from network server when it joins network

Can renew its lease on address in use

Allows reuse of addresses (only hold address while connected an "on") Support for mobile users who want to join network (more shortly) DHCP overview:

- host broadcasts "DHCP discover" msg [optional]
- DHCP server responds with "DHCP offer" msg [optional]
- host requests IP address: "DHCP request" msg
- DHCP server sends address: "DHCP ack" msg
- 14

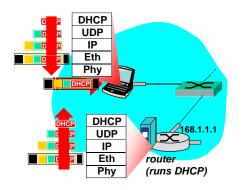
### **DHCP** client-server scenario



# DHCP: more than IP address

- DHCP can return more than just allocated IP address on subnet:
  - address of first-hop router for client
  - name and IP address of DNS sever
  - network mask (indicating network versus host portion of address)

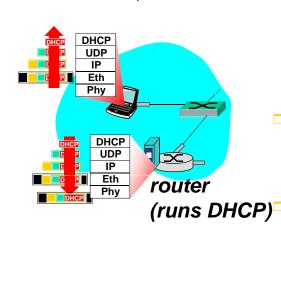
#### DHCP: example



- connecting laptop needs its IP address, addr of first-hop router, addr of DNS server: use DHCP
- DHCP request encapsulated in UDP, encapsulated in IP, encapsulated in 802.1 Ethernet
- Ethernet frame broadcast (dest: FFFFFFFFFFF) on LAN, received at router running DHCP server
- Ethernet demux'ed to IP demux'ed, UDP demux'ed to DHCP

18

#### DHCP: example



- DCP server formulates DHCP ACK containing client's IP address, IP address of first-hop router for client, name & IP address of DNS server
- encapsulation of DHCP server, frame forwarded to client, demux'ing up to DHCP at client

client now knows its IP address, name and IP address of DSN server, IP address of its first-hop router

### 21-2 ICMP

The IP protocol has no error-reporting or errorcorrecting mechanism. The IP protocol also lacks a mechanism for host and management queries. The Internet Control Message Protocol (ICMP) has been designed to compensate for the above two deficiencies. It is a companion to the IP protocol.

#### **Topics discussed in this section:**

Types of Messages Message Format Error Reporting and Query Debugging Tools

### ICMP: Internet Control Message Protocol

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- used by hosts & routers to communicate network-level information
  - error reporting: unreachable host, network, port, protocol
  - echo request/reply (used by ping)
- network-layer "above" IP:
  - ICMP msgs carried in IP datagrams
- ICMP message: type, code plus first 8 bytes of IP datagram causing error

- Type Code description
  - 0 echo reply (ping)
  - 0 dest. network unreachable 1
  - dest host unreachable 2
    - dest protocol unreachable
    - dest port unreachable
    - dest network unknown
    - dest host unknown
  - 0 source quench (congestion control - not used)
  - 0 echo request (ping) 0 route advertisement
  - 0 router discovery
  - 0 TTL expired
  - 0 bad IP header

## Traceroute and ICMP

- Source sends series of UDP segments to dest
  - First has TTL =1
  - Second has TTL=2, etc.
  - Unlikely port number
- When nth datagram arrives to nth router:
  - Router discards datagram .
  - And sends to source an ICMP message (type 11, code 0)
  - Message includes name of router& IP address

- When ICMP message arrives, source calculates RTT
- Traceroute does this 3 times

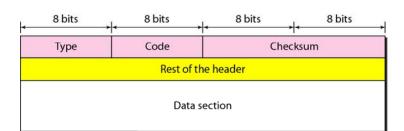
#### Stopping criterion

- UDP segment eventually arrives at destination host
- Destination returns ICMP "host unreachable" packet (type 3, code 3)
- When source gets this ICMP, stops.

21



#### Figure 21.8 General format of ICMP messages



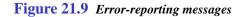


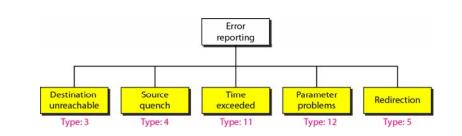
• Error-reporting message: reports problems that a host or a router may encounter when it processes an IP packet.

• Query message: helps a host or network manager get specific information form a router or another host

#### Note

ICMP always reports error messages to the original source.





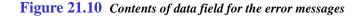
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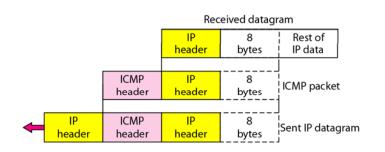
Note

26

### Important points about ICMP error messages:

- No ICMP error message will be generated in response to a datagram carrying an ICMP error message.
- No ICMP error message will be generated for a fragmented datagram that is not the first fragment.
- ❑ No ICMP error message will be generated for a datagram having a multicast address.
- No ICMP error message will be generated for a datagram having a special address such as 127.0.0.0 or 0.0.0.0.





#### Figure 21.11 Redirection concept

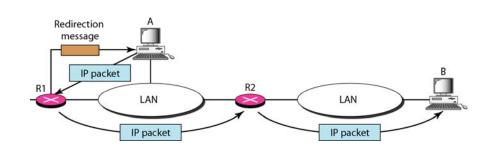
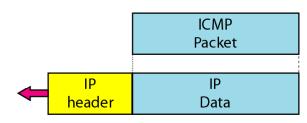
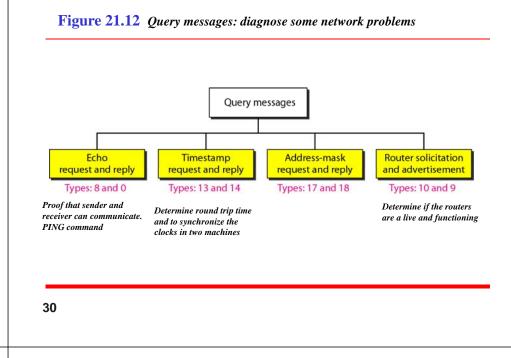


Figure 21.13 Encapsulation of ICMP query messages





#### Example 21.3

We use the ping program to test the server fhda.edu. The result is shown on the next slide. The ping program sends messages with sequence numbers starting from 0. For each probe it gives us the RTT time. The TTL (time to live) field in the IP datagram that encapsulates an ICMP message has been set to 62. At the beginning, ping defines the number of data bytes as 56 and the total number of bytes as 84. It is obvious that if we add 8 bytes of ICMP header and 20 bytes of IP header to 56, the result is 84. However, note that in each probe ping defines the number of bytes as 64. This is the total number of bytes in the ICMP packet (56 + 8).

### Example 21.3 (continued)

#### \$ ping fhda.edu

#### PING fhda.edu (153.18.8.1) 56 (84) bytes of data.

64 bytes from tiptoe.fhda.edu (153.18.8.1): icmp_seq=0	ttl=62	time=1.91 ms
64 bytes from tiptoe.fhda.edu (153.18.8.1): icmp_seq=1	ttl=62	time=2.04 ms
64 bytes from tiptoe.fhda.edu (153.18.8.1): icmp_seq=2	ttl=62	time=1.90 ms
64 bytes from tiptoe.fhda.edu (153.18.8.1): icmp_seq=3	ttl=62	time=1.97 ms
64 bytes from tiptoe.fhda.edu (153.18.8.1): icmp_seq=4	ttl=62	time=1.93 ms
64 bytes from tiptoe.fhda.edu (153.18.8.1): icmp_seq=5	ttl=62	time=2.00 ms
64 bytes from tiptoe.fhda.edu (153.18.8.1): icmp_seq=6	ttl=62	time=1.94 ms
64 bytes from tiptoe.fhda.edu (153.18.8.1): icmp_seq=7	ttl=62	time=1.94 ms
64 bytes from tiptoe.fhda.edu (153.18.8.1): icmp_seq=8	ttl=62	time=1.97 ms
64 bytes from tiptoe.fhda.edu (153.18.8.1): icmp_seq=9	ttl=62	time=1.89 ms
64 bytes from tiptoe.fhda.edu (153.18.8.1): icmp_seq=10	ttl=62	time=1.98 ms

#### --- fhda.edu ping statistics ---

11 packets transmitted, 11 received, 0% packet loss, time 10103ms rtt min/avg/max = 1.899/1.955/2.041 ms

35

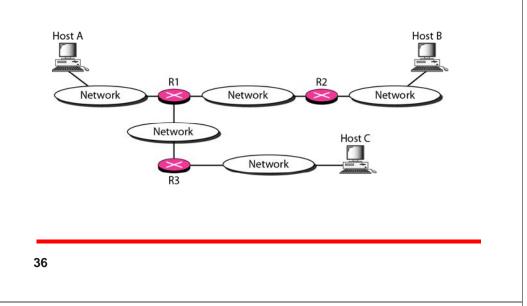
### Example 21.4

We use the traceroute program to find the route from the computer voyager.deanza.edu to the server fhda.edu. The following shows the result:

\$ traceroute fhda.edu				
traceroute to fhda.edu (153.18.8.1), 30 hops max, 38 byte packets				
1 Dcore.fhda.edu	(153.18.31.254)	0.995 ms	0.899 ms	0.878 ms
2 Dbackup.fhda.edu	(153.18.251.4)	1.039 ms	1.064 ms	1.083 ms
3 tiptoe.fhda.edu	(153.18.8.1)	1.797 ms	1.642 ms	1.757 ms

The unnumbered line after the command shows that the destination is 153.18.8.1. The packet contains 38 bytes: 20 bytes of IP header, 8 bytes of UDP header, and 10 bytes of application data. The application data are used by traceroute to keep track of the packets.

Figure 21.15 The traceroute program operation - tracert



### **Example 21.4 (continued)**

The first line shows the first router visited. The router is named Dcore.fhda.edu with IP address 153.18.31.254. The first round-trip time was 0.995 ms, the second was 0.899 ms, and the third was 0.878 ms. The second line shows the second router visited. The router is named Dbackup.fhda.edu with IP address 153.18.251.4. The three round-trip times are also shown. The third line shows the destination host. We know that this is the destination host because there are no more lines. The destination host is the server fhda.edu, but it is named tiptoe.fhda.edu with the IP address 153.18.8.1. The three round-trip times are also shown.

### Example 21.5

In this example, we trace a longer route, the route to xerox.com (see next slide). Here there are 17 hops between source and destination. Note that some roundtrip times look unusual. It could be that a router was too busy to process the packet immediately.



### 21-3 IGMP

The IP protocol can be involved in two types of communication: unicasting and multicasting. The Internet Group Management Protocol (IGMP) is one of the necessary, but not sufficient, protocols that is involved in multicasting. IGMP is a companion to the IP protocol.

#### **Topics discussed in this section:**

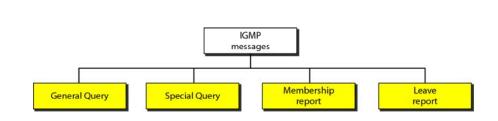
Group Management IGMP Messages and IGMP Operation Encapsulation Netstat Utility

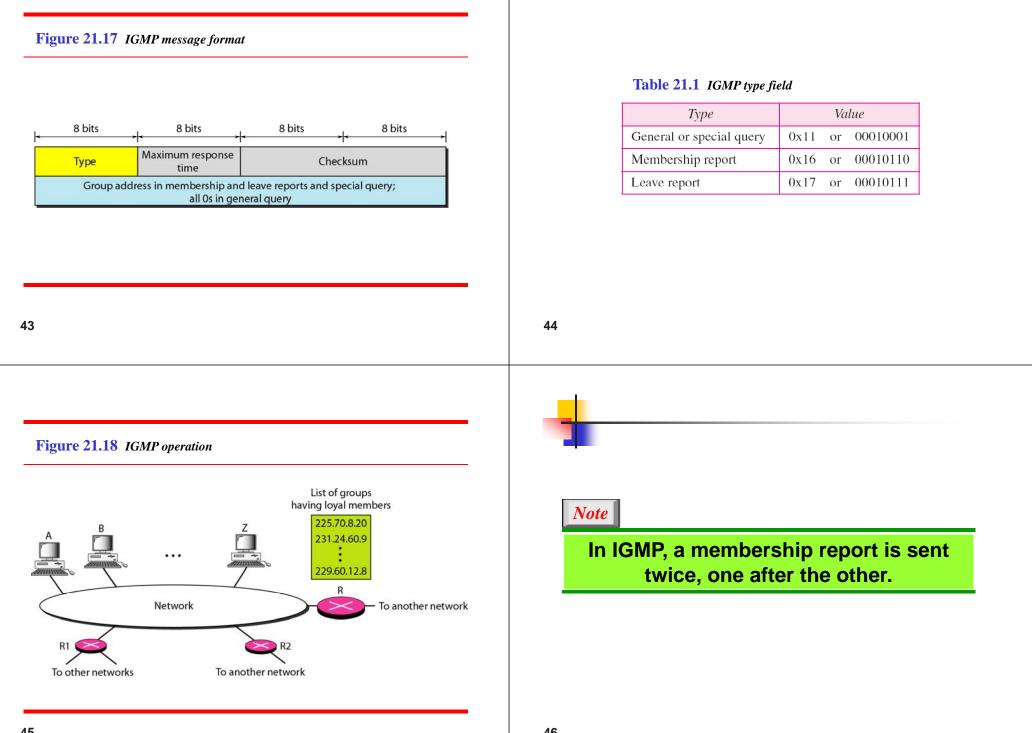
# **Example 21.5 (continued)**

\$ traceroute xerox.co				
traceroute to xerox.com (13.1.64.93), 30 hops max, 38 byte packets				
1 Dcore.fhda.edu	(153.18.31.254)	$0.622 \mathrm{ms}$	0.891 ms	$0.875 \ {\rm ms}$
2 Ddmz.fhda.edu	(153.18.251.40)	2.132 ms	2.266 ms	2.094 ms
3 Cinic.fhda.edu	(153.18.253.126)	2.110 ms	2.145 ms	1.763 ms
4 cenic.net	(137.164.32.140)	3.069 ms	2.875 ms	2.930 m
5 cenic.net	(137.164.22.31)	4.205 ms	4.870 ms	4.197 m
14 snfc21.pbi.net	(151.164.191.49)	7.656 ms	7.129 ms	6.866 m
15 sbcglobal.net	(151.164.243.58)	7.844 ms	7.545 ms	7.353 m
16 pacbell.net	(209.232.138.114)	9.857 ms	9.535 ms	9.603 m
17 209.233.48.223	(209.233.48.223)	10.634 ms	10.771 ms	10.592 m
18 alpha.Xerox.COM	(13.1.64.93)	11.172 ms	11.048 ms	10.922 m

40

**Figure 21.16** *IGMP message types* 







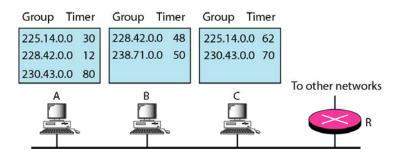
The general query message does not define a particular group.



Imagine there are three hosts in a network, as shown in Figure 21.19. A query message was received at time 0; the random delay time (in tenths of seconds) for each group is shown next to the group address. Show the sequence of report messages.

48

Figure 21.19 Example 21.6





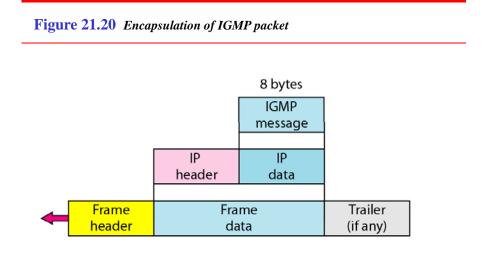
#### **Solution**

The events occur in this sequence:

a. Time 12: The timer for 228.42.0.0 in host A expires, and a membership report is sent, which is received by the router and every host including host B which cancels its timer for 228.42.0.0.

### Example 21.6 (continued)

- b. Time 30: The timer for 225.14.0.0 in host A expires, and a membership report is sent which is received by the router and every host including host C which cancels its timer for 225.14.0.0.
- c. Time 50: The timer for 238.71.0.0 in host B expires, and a membership report is sent, which is received by the router and every host.
- d. Time 70: The timer for 230.43.0.0 in host C expires, and a membership report is sent, which is received by the router and every host including host A which cancels its timer for 230.43.0.0.



52

### Note

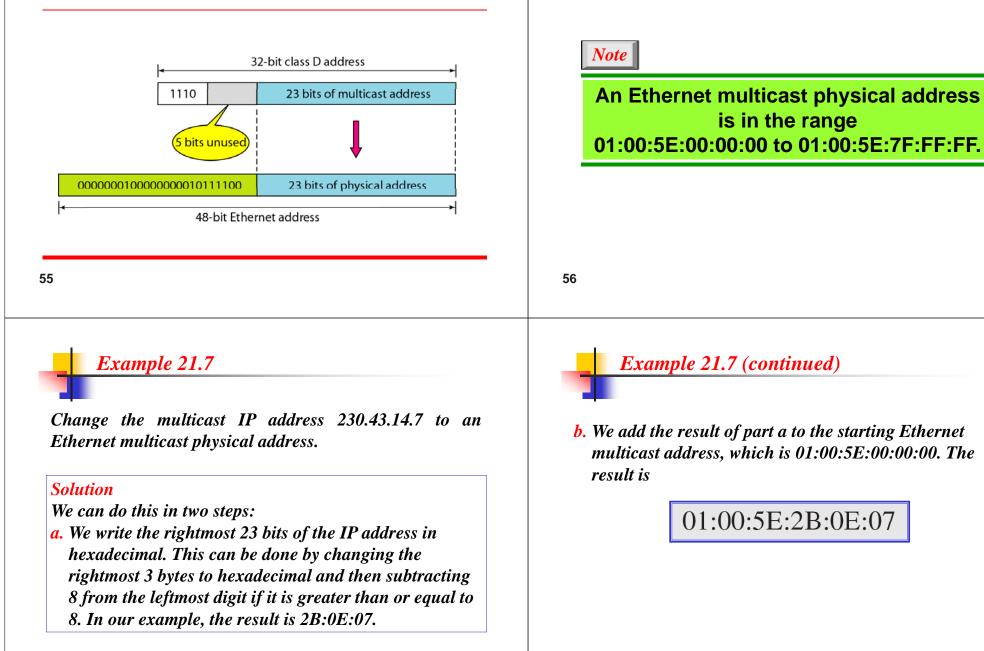
51

The IP packet that carries an IGMP packet has a value of 1 in its TTL field.

#### Table 21.2 Destination IP addresses

Туре	IP Destination Address
Query	224.0.0.1 All systems on this subnet
Membership report	The multicast address of the group
Leave report	224.0.0.2 All routers on this subnet





is in the range

Example 21.8

Change the multicast IP address 238.212.24.9 to an Ethernet multicast address.

#### **Solution**

- a. The rightmost 3 bytes in hexadecimal is D4:18:09. We need to subtract 8 from the leftmost digit, resulting in 54:18:09.
- **b.** We add the result of part a to the Ethernet multicast starting address. The result is

# 01:00:5E:54:18:09

59

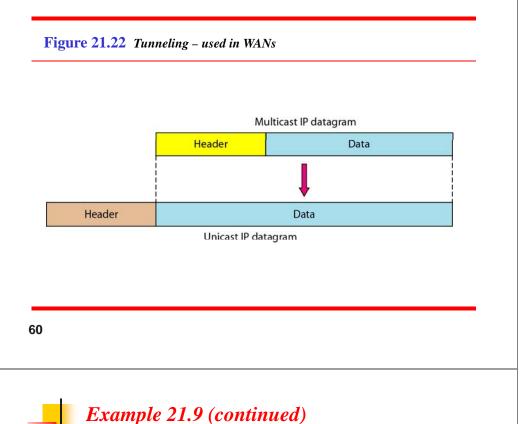
61

### Example 21.9

Netstat is used to find the multicast addresses supported by an interface

We use netstat (see next slide) with three options: -n, -r, and -a. The -n option gives the numeric versions of IP addresses, the -r option gives the routing table, and the -a option gives all addresses (unicast and multicast). Note that we show only the fields relative to our discussion. "Gateway" defines the router, "Iface" defines the interface.

Note that the multicast address is shown in color. Any packet with a multicast address from 224.0.0.0 to 239.255.255.255 is masked and delivered to the Ethernet interface.



Kernel IP routing table				
Destination	Gateway	Mask	Flags	Iface
153.18.16.0	0.0.0	255.255.240.0	U	eth0
169.254.0.0	0.0.0	255.255.0.0	U	eth0
127.0.0.0	0.0.0	255.0.0.0	U	lo
224.0.0.0	0.0.0	224.0.0.0	U	eth0
0.0.0.0	153.18.31.254	0.0.0.0	UG	eth0

#### 21-4 ICMPv6

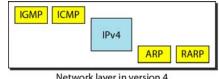
We discussed IPv6 in Chapter 20. Another protocol that has been modified in version 6 of the TCP/IP protocol suite is ICMP (ICMPv6). This new version follows the same strategy and purposes of version 4.

#### **Topics discussed in this section:**

**Error Reporting** Query

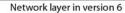
63

Figure 21.23 Comparison of network layers in version 4 and version 6





Network layer in version 4



64

#### Table 21.3 Comparison of error-reporting messages in ICMPv4 and ICMPv6

Type of Message	Version 4	Version 6
Destination unreachable	Yes	Yes
Source quench	Yes	No
Packet too big	No	Yes
Time exceeded	Yes	Yes
Parameter problem	Yes	Yes
Redirection	Yes	Yes

#### Table 21.4 Comparison of query messages in ICMPv4 and ICMPv6

Type of Message	Version 4	Version 6
Echo request and reply	Yes	Yes
Timestamp request and reply	Yes	No
Address-mask request and reply	Yes	No
Router solicitation and advertisement	Yes	Yes
Neighbor solicitation and advertisement	ARP	Yes
Group membership	IGMP	Yes